

Second Assignment

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- Interplay between a data structure, the append-only log, and the broadcasting communication abstraction.
- **Context:** Secure Scuttlebutt (SSB) project
 - <u>https://ssbc.github.io/scuttlebutt-protocol-guide/</u>
 - <u>https://handbook.scuttlebutt.nz/</u>
 - <u>https://github.com/ssbc/ssb-server</u>
- An overlay network that is tailored for decentralized social applications
- Networking with arbitrary data packets could be replaced by networking with coherent data structures.
- What are suitable data structures?
- What would corresponding networking primitives look like?

Middleware for Social Applications

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- The infrastructure is provided by participants.
- They are developed by open-source communities.
- Scuttlebutt has pioneered the use of replicated authenticated singlewriter append-only logs:
 - Chains of ordered immutable events specific to each participant.
 - Replicated by gossip algorithms that are driven by social signals.
- Two gossip models that can be used for replication:
 - **Open model** that works best in small and trusted groups.
 - Transitive-interest model that scales to thousands of participants.

Middlewares for Social Applications



- Minimal infrastructure requirements beyond participants' devices.
- <u>Hypercore</u>: to share large scientific datasets between research teams.
- Secure-Scuttlebut (SSB) is actively used for social applications
 - Blogging and code development.
 - Active open-source mobile clients.
- Both projects pioneered the use of replicated authenticated singlewriter append-only logs.
 - Chains of ordered immutable events, similar to blockchains.
 - Specific to each participant and without the need for global consensus.
- SSB replicates logs using a gossip protocol
 - By only transmitting the latest missing events between pairs of replicas.

SSB vs other Gossip Algorithms

- SSB has less parameters to tune
 - **No temporary** buffer size;
 - **Logs are persisted** because storage is abundant and cheap;
 - No retransmission, events are exchanged over reliable channels when found missing from meta-data;
 - No fan-out number of neighbours to contact
- The result is a simple yet practical messaging middleware.





- SSB organises events in logs, that are replicated between stores.
- Events may represent a user action or the result of processing operations.
- Logs organise events in a data structure for efficient replication and integrity.
- Stores hold many replicated logs either in the storage of an active process, or passively on a storage device.
- Stores are connected over reliable communication channels
 - E.g., TCP connections over the Internet or USB connections on a local machine.
- Many data formats and communications protocols can be used.
 - We can abstract from this decision.





- *id:* is the publicKey of the creator
- previous: is the hash of the previous event in the log including the signature. Null if none
- **index:** (sequence number) is the position of the event in the log
- **content**: is defined by applications
- signature: is the cryptographic signature of id, previous, index, and content, obtained with the privateKey that corresponds to the publicKey.
- All fields except content represent the meta-data about events.
- They enable efficient replication and integrity of the event chain





- A log is an identifier id associated to a sequence of events,
 - An empty log is id: [] and a log with two events is
 - id: [(id,null,0,..,.),(id,hash0,1,..,.)] with the following operations.

Table 1. Authenticated single-writer log operations

```
log \leftarrow create(publicKey)

log \leftarrow log.append(content, privateKey)

log \leftarrow log.update(events)

events \leftarrow log.get(start, end)

index \leftarrow log.last

id \leftarrow log.id
```

create a log from a public key

extend the log with a new event created locally (as owner) from *content* extend the log with the subset of compatible *events* previously created remotely get the set of events with index included between *start* and *end* (can be the same) get the index of the last event stored locally get the id (publicKey) of the log

- All operations are implemented to ensure the log is,
 - Secure
 - Monotonic
 - Linear (total order)
 - Single Writer
 - Connected



- A **Store** is a set {log1,...} that represents the logs stored locally.
- Adding or removing logs
 - Direct user actions, or
 - Indirectly the result of operations triggered by new events added to logs.
- A frontier is a set {(logi.id, logj.last),..} that represents the latest known indexes about logs in a store.
- The *difference between two frontiers* represents the new events in one store that are not yet in the other.



- Two frontiers may differ in the logs they contain because two stores may contain only partially overlapping sets of logs
- Stores and Frontiers are linked to events and logs by the following operations

$store \leftarrow store.add(log)$	add a log to the store
$store \leftarrow store.remove(id)$	remove the log with id
	from the store
$log \leftarrow store.get(id)$	get the log with <i>id</i> from
	the store (if present)
ids ← store.ids	get the set of ids of the
	logs in the store
$frontier \leftarrow store.frontier(ids)$	get the current <i>frontier</i> of
	the store only for <i>ids</i>
$events \leftarrow store.since(frontier)$	get the set of <i>events</i> that
	happened after <i>frontier</i>
$store \leftarrow store.update(events)$	update the logs in store
	with events



- Two stores in different locations may diverge because new events have been added locally.
- Updating is the process of propagating the new events in one store not present in the other.

Algorithm 1 *Update*(*store*, *store'*): update *store* and *store'* with missing new events present in the other.

- 1: frontier ← store.frontier(store.ids)
- 2: frontier' ← store'.frontier(store.ids)
- 3: $news \leftarrow store.since(frontier')$
- 4: $news' \leftarrow store'.since(frontier)$
- 5: $store \leftarrow store.update(news')$
- 6: $store' \leftarrow store'.update(news)$



- SSB uses Ed25519 key pairs both for signature and encryption.
- Access to content is restricted by encrypting it using the publicKey of recipient(s) using a private-box, a scheme that hides the recipients, their number, and the content.





• The goal of this model is to maximize the diffusion of events by replicating them in all participants.

1: l	оор
2:	Randomly pick store and store' from Participants
3:	for <i>id</i> in <i>store'.ids</i> – <i>store.ids</i> do <i>store.add</i> (<i>create</i> (<i>id '</i>))
4:	for <i>id</i> in <i>store.ids</i> – <i>store</i> '. <i>ids</i> do <i>store</i> '. <i>add</i> (<i>create</i> (<i>id</i>))
5:	Update(store, store')
6: e	nd loop



- Conversations between many participants, are spread over multiple logs.
- To ensure participants receive all messages, the logs in the transitive graph of interests should therefore be replicated.
- Additionally, this may serve to discover new interesting people to follow
 - Ex: Friend may introduce a friend to another one.
- The logs participating in the gossip algorithm propagate the interest/disinterest information.

Algorithm 3 Transitive-Interest Gossip \triangleright Participants abbreviated P and store abbreviated st.1: loop2: Randomly pick (id,st) and (id,st') from P3: for f in followed(st,id) – st.ids do st.add(create(f))4: for b in blocked(st,id) \cap st.ids do st.remove(b)5: ...6: Update(st, st')7: end loop

Second Assignment



- Deadline: Monday 4, January 2021 6pm
- Reference paper available on my website

- 1) <u>PDF document</u> reporting all the computational analysis of the implemented algorithm using the Simulator and a small description of the model designed (i.e., a tutorial to execute the model)
- 2) <u>A GitHub Repo</u> containing :
 - i. a README file that describes the members of the project and a summary of the implemented algorithm.
 - ii. Source code of the simulation.
 - iii. A JAR file of the model installer.